- 1. Prove or disprove or show that the statement is equivalent to an open question:
  - (a)  $\mathsf{DTIME}(2^n) \subsetneq \mathsf{NTIME}(2^{2n})$ .
  - (b)  $P \neq NP$  or  $NP \neq EXP$ .
  - (c) There exists a k > 0 such that NP  $\subseteq$  DTIME $(n^k)$ .
  - (d) For any language  $L \in \mathsf{NP} \cap \mathsf{coNP}, \mathsf{NP}^L = \mathsf{NP}$ .
- 2. (a) Let  $\Sigma_2$ SAT denote the following decision problem: given a quantified formula  $\psi$  of the form  $\psi = \exists x_1, \ldots, x_n \ \forall y_1, \ldots, y_n$ .  $\phi(x_1, \ldots, x_n, y_1, \ldots, y_n)$ , where  $\phi$  is a CNF formula, decide whether  $\psi$  is true. Prove that if P = NP then  $\Sigma_2$ SAT  $\in P$ .
  - (b) For  $k_1, k_2 \in \mathbb{N}$  define the problem  $(k_1, k_2)$ -Coloring Extension (in short,  $(k_1, k_2)$ -CE) as follows: given a graph *G* and a set of vertices *S*, decide whether any  $k_1$ -coloring of *S* can be extended to a  $k_2$ -coloring of *G*. Show that (2,3)-CE  $\in \Pi_2^p$  and that (2,2)-CE  $\in$  coNP.
- 3. (a) Prove that if  $NTIME(n) \subseteq DTIME(n^{10})$  then P = NP. Hint: First use a padding argument to show that for any  $k \ge 1$ ,  $NTIME(n^k) \subseteq DTIME(n^{10k})$ .
  - (b) Prove that if every unary NP-language is in P then EXP = NEXP, and conclude that if EXP ≠ NEXP then there exists a language L ∈ NP \ P that is not NP-complete. Remark: It is known that there exists a language L ∈ NP \ P that is not NP-complete assuming the weaker assumption P ≠ NP (Ladner's Theorem).
- 4. We define the class  $\mathbf{S}_2^p$  as the set of all languages *L* for which there exist a polynomial-time Turing machine *M* and a polynomial *p* such that for all  $x \in \{0, 1\}^*$ ,

$$\begin{aligned} x \in L \Rightarrow \exists y \in \{0,1\}^{p(|x|)} \ \forall z \in \{0,1\}^{p(|x|)}. \ M(x,y,z) = 1 \\ x \notin L \Rightarrow \exists z \in \{0,1\}^{p(|x|)} \ \forall y \in \{0,1\}^{p(|x|)}. \ M(x,y,z) = 0 \end{aligned}$$

- (a) Is  $\mathbf{S}_2^p$  closed under complement?
- (b) Prove that  $\mathbf{S}_2^p \subseteq \Sigma_2^p \cap \Pi_2^p$ .