<u>Computational Models - Exercise 3</u> 8/12/08

- 1.
- a. Write a Turing machine that computes the function f(w) = 101w. Give a complete formal description of the machine. $(\sum = \{0, 1\})$
- b. Write a Counter machine that computes the following function: $f(n) = \lfloor n/2 \rfloor$. You may use any number of registers. Give a complete formal description of the machine (i.e., the number of registers and a finite sequence of instructions as explained in recitation and class, see definition 1 in chapter 3.2 in L1.pdf in the site of the course).
- 2.
- a. Write a Turing machine that decides the language $L_2 = \{I^k | k=3^n, n \in \mathbb{N}\}$ ($\sum = \{I\}$). Don't give a formal description, only describe the way the machine works (in the spirit of what was done in recitation).
- b. Write a counter machine that computes the function $f(n) = n^2$. Don't give a formal description, only describe the way the machine works (in the spirit of what was done in recitation).
- 3. We define the following variations of Turing machines:
 - a. Machines of the type TM' that are similar to regular TMs, except that |Q| < 2008 and $|\Gamma| < 2008$. In addition the head of the tape can move up to 2008 cells to the right or left at one step. For example: $\delta(q_5, a) = q_{17}, b, R28$ means that if the tape is at state q_5 and reads the character *a*, it shall move to state q_{17} write *b* on the tape and move 28 cells to the right. Let L(TM') be the languages accepted by TM' machines. Decide and prove whether $L(TM') \subset RE, RE \subset L(TM')$ or L(TM') = RE.
 - b. Machines of the type TM'' that are similar to regular TMs, except that they can move to the right or left any number of cells. Decide and prove whether RE=L(TM'') or $RE\subset L(TM'')$.
- 4. For the following languages determine whether they belong to *R*, *RE/R*, *co-RE/R*, or none of the above, and prove correctness:
 - a. Input: A Turing machine *M* and a natural number *K*. Question: Is there an input for which *M* halts in less than *K* steps.
 - b. Input: A Turing machine *M* and a string *x*. Question: Does M write '1' on the tape during execution on *x*.
 - c. Input: A Turing machine MQuestion: is M a Turing machine such that for any input x it does not reach the |x|+1000 cell. Hint: Use configurations and computational histories.
- 5. The accepting language $Accept = \{ < p, x > | p \text{ is a program that accepts } x \}$ is in RE\R. Prove that $L \in RE$ iff $L \leq_m Accept$.
- 6. Prove or disprove:
 - a. The class *RE* is closed under union and intersection
 - b. The class *co-RE* is closed under union and intersection

- c. If the language A is not decidable and $A \leq_m A^c$ (The complement of A), then $A \notin RE$ and $A \notin co-RE$
- 7. Let $L_1, L_2 \subseteq \{0, 1\}^*$ be *RE* languages such that $L_1 \cup L_2 = \{0, 1\}^*$ and $L_1 \cap L_2 \neq \phi$. Show that $L_1 \leq_m L_1 \cap L_2$.